N-LANE BRIDGE PERFORMANCE ANTIPATTERN ANALYSIS USING SYSTEM-LEVEL EXECUTION TRACING

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THE PROBLEM

- Multithreading can add latency to an application when used incorrectly
- The cause of this latency can be difficult to diagnose
 - Many possible causes: lock contention, resource pools, CPU preemption, etc.
 - Non-deterministic nature means that it is not trivial to follow the program's execution



- Surface indication of a deeper problem within the system
- Closely related to static analysis: the bad smell must be directly identifiable in source code
- Common examples: duplicate code, feature envy
- Previously, we worked on detecting runtime smells using execution tracing
 - Examples: CPU Hog, Priority Inversion
 - https://github.com/riley-v/runtime-badsmell-trace-metrics

- Common problem bundled with detection methods and common solutions
- Used to better describe issues in a piece of software
- Antipattern analysis is performed during the testing phase of an application, before delivery
- One Lane Bridge: a performance antipattern
 - Only one, or a few, threads can execute at once at this place

BAD SMELLS

ANTIPATTERNS



EXISTING CHALLENGES

One Lane Bridge Performance Antipattern

- Bundles a common multithreading bottleneck problem with detection strategies and refactoring solutions
- Two issues/gaps in existing research:
 - 1. Do not consider bottlenecks in active resources
 - ► Active resource: performs a physical action in the real world (eg. CPU)
 - vs. passive resource: exists only virtually (eg. mutex)
 - 2. Use imprecise metrics
 - Example: overall CPU usage is used to denote application's CPU access



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OUR SOLUTION

- Introduce N-Lane Bridge: new category of One Lane Bridge
 - Extends OLB from passive resources to active resources
 - Defines a method to distinguish an NLB in the target application from issues due to an external application
- Introduce detection method using system-level execution tracing to be performed during the testing phase
 - Allows for the gathering of more precise metrics
 - Eliminates the need for manual instrumentation of source code with tracepoints



WHAT IS THE N-LANE BRIDGE ANTIPATTERN?

Definition:

- A performance bottleneck due to the target application's use of an active resource.
- Cases where the bottleneck is due to an external application's use of an active resource are not considered NLB's.

Probable Causes:

- Incorrect configuration of the software.
- Implementation of the software did not account for the system it would be running on.



WHY FOCUS ON SYSTEM-LEVEL TRACING?

Benefits for Multithreading:

- Concurrency is difficult to predict. Dynamic analysis eliminates the false alarms from abstraction found in static analysis.
- System-level tracing provides the wide visibility of executing threads, processes, and resources needed to accurately analyze multithreaded applications.

Benefits for Ease of Use:

The tracepoints are predefined (eg. Linux kernel) so no extra effort or domain knowledge is needed.



DETECTION STRATEGY





Log In

Username		
Password		
	Login	

TRACE COLLECTION

Stop Conditions:

- The response time has increased 10x its original number
- 2. The application no longer supports adding more users

Required Tracepoints:

- Syscalls: reveal the reason for a blocked thread
 - Interrupt syscalls: irq_handler, softirq, hrtimer_expire
- Request delimiters: reveal the response time for a request
- sched_switch: indicate when a thread occupying a CPU is switched for the new thread
- sched_wakeup: indicate when a previously blocked thread becomes runnable



1ThreadTest.jmx (/home/riley/Documents/apache-jmeter-5.4.1/bin/1ThreadTest.jmx) - Apache JMeter (5.4.1)

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	🖋 Login

🎿 Backend Listener

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	Basic Advanced				
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	HTTP Request				
	POST				Content encoding:
	Redirect Automatically 🗹 Follow Redirects 🗹 Use KeepAlive	Use multipart/form-data 📃 Browser-compatible header			
	Parameters Body Data Files Upload				
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	pwd	password		text/plain	
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METRIC EXTRACTION

Metrics Needed:

- Response times of requests
- Critical blocking times for threads handling the requests, as well as their causes

Critical Path:

- Longest series of steps before completion
- Can be used to find critical blocking times and their causes



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RESPONSE TIME ANALYSIS

Motivation:

 Both categories of One Lane Bridge cause increasing response time latency as more users stress the system

Analysis Formula: $\exists y \forall z (RT_z < RT_{z+1}) \rightarrow P$ Where:

- RT_z is the average response time for any execution z
- Any execution z has less users than execution z+1
- P holds that response times increase after execution y

Comparison of response times is done using a two-tailed t-test.



Open	(F)
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1 Start, 289

2 Start, 821, 1718, 1719, 1735, 1726, 1736

3 Start, 2130, 2132, 2148, 2184, 2180, 2169, 2176, 2176, 2159, 2126, 2226

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4 Start, 1756, 2695, 2740, 2760, 2862, 2897, 2955, 2959, 3000, 3015, 3023, 3039, 3055, 3067, 3082, 3087 **5** Start, 2971, 3137, 3549, 3604, 3636, 3838, 3850, 3864, 3908, 3951, 3960, 3950, 3947, 3961, 3964, 3975, 3960, 3976, 3982, 4003, 4005

LATENCY CAUSE ANALYSIS

Average Blocking Time Formula: $AveBlockTime(Res_x) = \sum_{i=y}^{n-1} (BTx_{i+1} - BT_{x_i})$ Where:

- Res_x is a resource which blocks the target threads
- BT_{xi} is the blocking time for resource x in execution1

Analysis Formula:

 $\exists Res_x \forall Res_z (AveBlockTime(Res_x) > AveBlockTime(Res_z)) \rightarrow Q$

Where:

 Q holds that Res_x is the resource which contributes the most to the latency





ACTIVE RESOURCE CONGESTION ANALYSIS Check which threads are using the congested resource during the analysis period:

- If over a specific threshold (e.g. 50%) of time, the target application holds the resource: N-Lane Bridge
- 2. Otherwise: No detected problem in the target application



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Address translation formula: *TracepointAddress – BaseAddress = ELFSymbolTableEntry*

CALL STACK ANALYSIS

0000000000004040	В	lock
000000000000134a	Т	main
	U	printf@@GLIBC_2.2.5
	U	pthread_create@@GLIBC_2.2.5
	U	pthread_join@@GLIBC_2.2.5
	U	pthread_mutex_destroy@@GLIBC_2.2.5
	U	pthread_mutex_init@@GLIBC_2.2.5
	U	pthread_mutex_lock@@GLIBC_2.2.5
	U	pthread_mutex_unlock@@GLIBC_2.2.5
	U	puts@@GLIBC_2.2.5
0000000000001200	t	register_tm_clones
00000000000011a0	Т	_start
	U	strerror@@GLIBC_2.2.5
0000000000004070	В	tid
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0000000000001289	Т	trythis

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No	data											

We have written a paper on this subject and submitted it to the 2022 ICPC conference.

Future Work:

- Improvement: the metric extraction algorithm takes a significant amount of time. This is due to the number of critical paths that must be formed and examined individually. An improvement would be to form and examine the critical paths in batch jobs, reducing the number of times the trace must be iterated over.
- Addition: more performance antipatterns could be detected using systemlevel tracing. Antipatterns similar to the One Lane Bridge include Traffic Jam and The Ramp.

CONCLUSIONS AND FUTURE WORK

Email: rv18jq@brocku.ca

GitHub Repo: <u>https://github.com/riley-v/n-lane-bridge-antipattern-analysis</u>

THANK YOU

COST AWARE TRACING

AMIR HAGHSHENAS MICHEL DAGENAIS NASER EZZATI

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PROGRESS REPORT 2022

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CONTENT





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COST AWARE TRACING



- Tracing generates large amount of data in short time
- It is not always possible to trace and save everything
 - IoT devices
 - Embedded systems
- A solution is required to choose how to collect as much as possible
 - Specially in systems with limited resources
- The solution should consider the observer effect of tracing on program
 - Execution delay
 - Memory consumption

COST AWARE TRACING



- How to reduce the cost
 - Sampling
 - Based on cost
 - Keep the most important



ANALYZING TRACING OBSERVER EFFECT



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- Fischmeister and Lam (2010) [3]
- Introduced time-aware instrumentation
- Cost functions are execution time and instrumentation coverage
- Starts with source code analysis and naïve instrumentation
- Check execution time
 - If violated, use integer linear programming to reduce coverage to meet budget.

USING COST FUNCTION TO MODIFY TRACING



- Kashif and Arafa (2013) [4]
- Static instrumentation model called INSTEP
- Support up to four extra-functional properties
 - Instrumentation intent values
 - Code size, execution time and detection latency
- Each cost model is created using an automata
- provided insight into pruning the search space of instrumentation alternatives to find a feasible instrumentation solution.

USING COST FUNCTION TO MODIFY TRACING



- Arafa and Kashif (2013) [5]
- DIME: Time-aware dynamic binary instrumentation
- Keep two version of the program.
 - One is instrumented and the other one is not instrumented
- User should provide two elements
 - P: Instrumentation period
 - B: Instrumentation budget (B<P)
- During each period, program is instrumented only for B time unit and then switch to the not instrumented version.

USING COST FUNCTION TO MODIFY TRACING



PROJECT OBJECTIVE



Define a budget for tracing overhead and monitor the system in real time.



Define a model for cost and effectiveness of each enabled trace point



Optimize the cost and effectiveness of the trace using an optimization problem

CURRENT STATE OF PROJECT



- Define a cost function to analysis
 - Execution delay
- Measurement
 - Developing kernel modules to record trace points execution time
 - Recording kernel counters using Perf_events
- Visualization
 - Working on trace compass development environment
 - Working on Histogram section to show the execution delay instead of event count



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THANKYOU





Adaptive Tracing Based On Time Series Trend Detection

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Challenges with execution tracing

□ Huge size and lots of noise and irrelevant data.

• High overhead (observer effect).

□ To many instruments: a typical Linux kernel has over 300 instruments.

- Its not easy to predict where system issues may arise and only conveniently keep those related instruments enabled.
- Requires manual tuning of tracepoints and instruments to make adjustments, which can be tedious.
- Pre-determined set of instruments might miss out on important tracing data.



Tracing methods

□ Static tracing – works by using static instrumentation.

- Dynamic tracing tracepoints are injected into a running system.
- □ Adaptive tracing tracing instruments/tracepoints are controlled automatically by the framework itself to figure out points of interest.

What is Adaptive Tracing?



Some related works

System Execution Path Profiling[1], LogMine [2], Pythia[3]:

- Using Coefficient of Variance (CoV) or any standard deviation methods on grouped data is more biased towards shorter requests.
- Clustering methods would always prioritize high standard deviation groups, where as there could be underlying consistently worse performing offenders going undetected.
- Anomaly based adaptive tracing would trigger events which are not always necessary.
- ADRL[4], DRAVM for Cloud Computing Environment[5], Workload Prediction Using ARIMA for Cloud Applications' QoS[6]:
 - Unlike cloud computing resource allocation via reinforcement learning methods, they are not feasible in tracing due to the lack of options to verify the actions right away.
- Process Monitoring on Sequences of System Call Count Vectors[7], Adaptive Performance Anomaly Detection in Distributed Systems Using Online SVMs[8]:
 - Trend based anomaly detection can't adapt to increasing or decreasing of the baseline resource usage.

Research questions

> How do we decide to change the tracing config?

- > Should we adjust trace config whenever there is an anomaly/change?
- > What are the possible actions? How we are going to change events?
- > How do we evaluate the trace adjustments?



Proposed method

□ Architecture / Structure:

- > Tracing
 - ***** Kernel-level tracing with LTTng
- Metrics Monitoring Layer
 - cpu_utilization, disk_utilization, etc.
 - Total of 23 metrics.
- > Trend Detection Layer
 - **Time series ARIMA, EMA, DEMA, etc.**
 - Financial market trend prediction
- > Anomaly Monitoring Layer
 - Maintains anomaly score
 - Adaptive threshold
- > Action layer
 - ***** Controls tracing events



TAT method flowchart



□ Metrics monitoring layer:

- > LTTng: Trace data collection
- Babeltrace: Metrics data extraction from trace files.

□ Trend detection layer:

- > Sampling:
 - Take sample of fixed size dataset at random intervals.
 - Control frequency of intervals based on anomaly scores.
- Use prediction model (ARIMA) to predict X+1 step.
- If mean of sample vs prediction is above or below anomaly threshold, then flag it as anomaly.

Proposed method

□ Anomaly monitoring layer:

- anomalyScore = (betaVal * isAnomaly) + ((1 betaVal) * anomalyScore)
- Auto adjust betaVal to increase or decrease anomalyScore based on *frequency of the anomaly occurrence*.

□ Action layer:

- Group LTTng events based on the list of relevant monitoring metrics.
- List problematic metrics and enable related LTTng events.
- Monitor changes in anomaly list to disable back events which are not required.

Demo



Example analysis



Conclusion & future work

- Binning of metrics data so that all metrics don't have to be processed through the trend detection layer all the time.
- □ Implement better method for a dynamic betaValue for the anomaly score function.
- □ Needs better correlation groupings of events and metrics list.
- □ Implement better methods to solve conflicts/overlap between events list.
- **Rigorous testing**.

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The End

Questions?



Thank you!